android

Android: protecting the kernel

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- Android Security since 2014
- Focus on system hardening
- Software Engineer



First - some good news

- Most kernel bugs discussed are not directly reachable to untrusted code, due to Android's security model. Android Nougat further reduces the attack surface of the kernel.
- New kernel defenses address our biggest category of kernel bugs. (more on this later)





Agenda

- Kernel bugs in Android
 - Focus on biggest categories (we only have 45 minutes)

• Bugs by cause

- Mitigations memory protections
- o Gaps
- Bugs by reachability
 - Mitigations attack surface reduction
 - Gaps
- Future work

Kernel bugs have a long life. Fixing bugs is necessary but not sufficient!



Goal:

Use data on security vulnerabilities to prioritize mitigation development and adoption

About the data

- Includes kernel bugs January 2014 -> April 2016
- Includes low \rightarrow critical severity kernel bugs
- Moderate \rightarrow critical taken directly from public Nexus security bulletins
- Low severity bugs included because the definition of low has changed over time (many bugs previously listed as low considered moderate under new ratings)

"At the operating system level, the Android platform provides the security of the Linux kernel..." source.android.com



Security from the kernel

- Address space separation/process isolation
- unix permissions
- DAC capabilities
- SELinux
- seccomp

. . .

namespaces



Android security bugs by year



Why the rise in kernel bugs?

- Lockdown of userspace makes UID 0 significantly less useful.
- 2016 is the first year > 50% of devices in ecosystem have selinux in global enforcing.
- Kernel bugs payout more \$\$\$ (Android vulnerability rewards)



Where Android's kernel bugs are born



Data includes multiple vendors

Some vendor drivers are in upstream, others are out-of-tree

Data: Jan 2014 \rightarrow April 2016



Kernel defenses protect against in-tree AND out-of-tree vulnerabilities

What causes Android's kernel bugs?



What causes kernel bugs?







Core kernel

Vendor drivers

Data: Jan 2014 \rightarrow April 2016

Mitigations - missing/incorrect bounds check

- Hardened usercopy
 - Protect against incorrect bounds checking in copy_*_user()
- PAN emulation
 - Protect against kernel access to userspace bypassing hardened usercopy changes.

Landing in upstream kernel!



Mitigations - missing/incorrect bounds check

- Stack protector strong
 - protects against stack buffer overflows
- KASLR (arm64 android-4.4 kernel)
 - Makes code reuse attacks probabilistic
- PXN make userspace non-executable for the kernel
 - Protects against ret2user attacks
- RODATA mark kernel memory as read-only/no-execute
 - Makes code non-writeable, and data non-executable



Mitigations - null pointer dereference

- CONFIG_LSM_MMAP_MIN_ADDR
 - Make null pointer dereference unexploitable (just crash)
- PAN emulation also make null pointer dereference non-exploitable



Gaps - code review

Code quality of upstream is significantly better than device specific drivers

- What can be done to enforce better code quality?
 - Compiler changes e.g. integer overflow checking
 - Scripts e.g. checkpatch.pl
 - Runtime changes e.g. PAN enforce proper use of copy_*_user()
 - KASAN
 - Constification



Attack surface reduction

Remove or restrict access to entry points into the kernel

Attack surface reduction

Gate access to kernel provided developer tools in developer settings.

How are kernel bugs reached - driver/subsystem



Data: Jan 2014 \rightarrow April 2016



Bugs reachable by unprivileged apps

Fix all bugs, but prioritize mitigation development for bugs that are reachable by apps

More on this later...



Data: Jan 2014 \rightarrow April 2016



Case study: Wifi driver bugs

- Every app-reachable bug should have been protected by a CAPABLE(CAP_NET_ADMIN) check.
- Relying on developers to correctly implement in-code checks is risky.
- Better to have privileged behavior guarded by auditable security policy.
- Many wifi driver bugs were reachable via local unix sockets! Add strong policy around all socket types.



How are kernel bugs reached - syscall (before mitigations)





100% of perf vulns introduced in vendor customizations



bugs reachable by apps

Data: Jan 2014 \rightarrow April 2016

Mitigations - attack surface reduction

loctl command whitelisting in SELinux

- Wifi
 - Originally hundreds of ioctl commands \rightarrow 29 whitelisted safe network socket ioctls
 - Blocks access to all bugs without restricting legitimate access.
 - Unix sockets: wifi ioctls reachable by local unix sockets :(Hundreds \rightarrow 8 whitelisted unix socket ioctls
 - No ioctls allowed on other socket types including generic and netlink sockets
- GPU

- e.g. Shamu originally 36 -> 16 whitelisted commands
- loctl commands needed varies by device but < 50% needed seems consistent across KGSL drivers

Mitigations - attack surface reduction

- Restrict access to perf
 - Access to perf_event_open() is disabled by default.
 - Developers may re-enable access via debug shell
- Remove access to debugfs
 - All app access to debugfs removed in N
- Remove default access to /sys
 - App access to files in /sys must be whitelisted
- Seccomp required for all devices (minijail shoutout!)

Impact of mitigations

Because most bugs are driver specific, effectiveness of mitigations varies across devices. In general most previously reachable bugs were made unreachable

- Case study of bugs reachable by apps on Nexus 6 (Shamu)
 - 100% of wifi bugs blocked
 - \circ $\,$ 50% of GPU bugs blocked $\,$
 - 100% of debugfs bugs blocked
 - 100% of perf bugs blocked (by default)



Gaps - Attack surface reduction

- Need more/better controls over kernel feature accessibility.
 - Controls allow us to do what's best for both Linux developers and users of Linux based products.
- Argument inspection for seccomp



Future work

Kernel devs, we need more/better safety features (seat belts)!

Sometimes seat belts are inconvenient...



Those "other" categories - potential attack surface reduction

Architecture	syscalls provided by kernel	syscalls in bionic	reduction (%)
arm	364	204	44
arm64	271	198	27
x86	373	203	46
x86_64	326	199	39



Those "other" categories - Memory safety



Where do we go from here?

- Kernel self protection project get involved!
- Principle of Least Privilege
- Attack Surface Reduction
- Defense-in-depth
- Continue to find/fix bugs!





